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Minisymposium Orthogonal Polynomials and Special Functions

Dual bases and orthogonal polynomials

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Part I. Dual bases: construction and applications

Dual basis: definition

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2. $\left\langle \mathbf{b}_i, \mathbf{d}_j^{(n)} \right\rangle = \delta_{ij} \quad (0 \leq i, j \leq n)$,

where $\delta_{ij} = 0$ for $i \neq j$, and $\delta_{ii} = 1$

Dual bases: applications

- Representation. If $g \in \mathcal{B}_n$ then

$$g = \sum_{k=0}^n a_k b_k, \quad \text{where } a_k := \langle g, d_k^{(n)} \rangle \quad (0 \leq k \leq n)$$

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- Approximation. For a given function f , the element

$$\mathbf{p}^* := \sum_{k=0}^n c_k \mathbf{b}_k, \quad \text{where } c_k := \langle f, \mathbf{d}_k^{(n)} \rangle \quad (0 \leq k \leq n),$$

is the best approximation of f in the set \mathcal{B}_n in the following sense:

$$\|f - \mathbf{p}^*\| = \min_{\mathbf{p} \in \mathcal{B}_n} \|f - \mathbf{p}\| \quad (\|\cdot\| := \sqrt{\langle \cdot, \cdot \rangle})$$

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Remarks.

1. One can find the optimal element *omitting* the orthogonal basis of the space \mathcal{B}_n .
2. Skillful use of dual bases often results in less costly algorithms of solving many computational problems.

Dual bases: recent results and applications

- Dual polynomial bases: Goldman (1994).
- Dual Bernstein polynomials: Ciesielski (1987), Jüttler (1998), Lewanowicz and W. (2006–2012), Rababah and Al-Natour (2007), Chudy and W. (2018, 2021).
- Dual Bézier-Said-Wang basis: Zhang *et al.* (2010).
- Computing roots of polynomials: Barton and Jüttler (2007), Liu *et al.* (2009).
- Degree reduction or merging of Bézier curves and surfaces: Gospodarczyk, Lewanowicz, W. (2009–2017).
- Polynomial approximation of rational Bézier curves and surfaces: Lewanowicz, Keller, W. (2012–2017).
- Construction of dual bases: W. (2013, 2014).
- Dual basis functions in subspaces: Kersey (2013).
- Method of solving boundary value problems: Gospodarczyk, W. (2017).
- Numerical solving of fractional partial differential equations: Jani *et al.* (2017, 2018)

Construction of dual basis: first method

- Let us write the dual function $d_i^{(n)}$ as a linear combination of b_0, b_1, \dots, b_n :

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- One can easily prove that the i th column of the following matrix:

$$\begin{bmatrix} \langle \mathbf{b}_0, \mathbf{b}_0 \rangle & \langle \mathbf{b}_1, \mathbf{b}_0 \rangle & \dots & \langle \mathbf{b}_n, \mathbf{b}_0 \rangle \\ \langle \mathbf{b}_0, \mathbf{b}_1 \rangle & \langle \mathbf{b}_1, \mathbf{b}_1 \rangle & \dots & \langle \mathbf{b}_n, \mathbf{b}_1 \rangle \\ \vdots & \vdots & \vdots & \vdots \\ \langle \mathbf{b}_0, \mathbf{b}_i \rangle & \langle \mathbf{b}_1, \mathbf{b}_i \rangle & \dots & \langle \mathbf{b}_n, \mathbf{b}_i \rangle \\ \vdots & \vdots & \vdots & \vdots \\ \langle \mathbf{b}_0, \mathbf{b}_n \rangle & \langle \mathbf{b}_1, \mathbf{b}_n \rangle & \dots & \langle \mathbf{b}_n, \mathbf{b}_n \rangle \end{bmatrix}^{-1}$$

contains the coefficients

$$\left[\mathbf{a}_{i0}^{(n)}, \mathbf{a}_{i1}^{(n)}, \dots, \mathbf{a}_{in}^{(n)} \right]^T \quad (0 \leq i \leq n)$$

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- Cost: $O(n^3)$

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Theorem. Let α_{ki} be the coefficients in

$$P_k = \sum_{i=0}^n \alpha_{ki} b_i \quad (k = 0, 1, \dots, n)$$

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Theorem. Let α_{ki} be the coefficients in

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Then the dual basis functions have the following form:

$$d_i^{(n)} = \sum_{k=0}^n \alpha_{ki} P_k \quad (i = 0, 1, \dots, n)$$

Construction of dual basis: third method

Disadvantages of the methods:

- we have to invert a full matrix (the first method),
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This method does not require:

- the inverting of a full matrix,
- the knowledge of the orthonormal basis

Construction of dual basis: third method – idea

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for the basis \mathbf{B}_k wrt. $\langle \cdot, \cdot \rangle$

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- It is possible to find the dual basis

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for $\mathbf{B}_{k+1} := \mathbf{B}_k \cup \{\mathbf{b}_{k+1}\}$ wrt. the same inner product in a linear time

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- There are quite simple connections between elements of the dual bases \mathbf{D}_k and \mathbf{D}_{k+1}

Construction of dual basis: third method – remarks

- Using this method, we can construct all the dual bases

$$D_1, D_2, \dots, D_n$$

in the time $O(n^2)$

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Problem. Given a function f and $\varepsilon > 0$, find the least natural number n such that

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Part II. Dual Bernstein and orthogonal polynomials

Dual Bernstein bases: classical case

- Basis functions \Rightarrow classical Bernstein basis polynomials (Bernstein, 1912):

$$B_i^n(x) = \binom{n}{i} x^i (1-x)^{n-i} \quad (0 \leq i \leq n)$$

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$$\langle f, g \rangle_J := \int_0^1 (1-x)^\alpha x^\beta f(x) g(x) dx \quad (\alpha, \beta > -1)$$

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- More on dual Bernstein polynomials and their relations with shifted Jacobi and Hahn (or dual Hahn) polynomials in a few minutes...

Dual Bernstein bases: discrete case

- Basis functions \Rightarrow discrete Bernstein basis polynomials (Neamtu, 1991; Sablonnière, 1992):

$$b_i^n(x; \mathbf{N}) = \frac{1}{(-\mathbf{N})_n} \binom{n}{i} (-x)_i (x - \mathbf{N})_{n-i} \quad (0 \leq i \leq n \leq \mathbf{N}; \mathbf{N} \in \mathbb{N})$$

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- Inner product \Rightarrow Hahn polynomials $Q_k(x; \alpha, \beta, N)$:

$$\langle f, g \rangle_H := \sum_{x=0}^N \binom{\alpha + x}{x} \binom{\beta + N - x}{N - x} f(x) g(x) \quad (\alpha, \beta > -1)$$

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$$d_i^n(x; \alpha, \beta, N) = \sum_{k=0}^n V_{Nk}^{(\alpha, \beta)} Q_k(i; \alpha, \beta, n) Q_k(x; \alpha, \beta, N)$$

Dual Bernstein bases: q -case

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$$B_k^n(x; \omega|q) = \frac{1}{(\omega q; q)_n} \begin{bmatrix} n \\ k \end{bmatrix}_q x^k (\omega q/x; q)_k (x; q)_{n-k} \quad (0 \leq k \leq n)$$

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$$D_i^n(x; a, b, \omega|q) = \sum_{k=0}^n V_k^{(a,b,\omega)}(q) Q_k(q^{i-n}; a, b, n|q) P_k(x; a, b, \omega|q)$$

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- Bernstein basis polynomials of degree n

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$$D_0^n(x; \alpha, \beta), D_1^n(x; \alpha, \beta), \dots, D_n^n(x; \alpha, \beta) \in \Pi_n$$

defined so that

$$\langle D_i^n, B_j^n \rangle_J = \delta_{ij} \quad (i, j = 0, 1, \dots, n),$$

where

$$\langle f, g \rangle_J := \int_0^1 (1-x)^\alpha x^\beta f(x) g(x) dx \quad (\alpha, \beta > -1)$$

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- Shifted Jacobi polynomials $R_k^{(\alpha, \beta)}(x)$ are orthogonal wrt the inner product $\langle f, g \rangle_J$, i.e.,

$$\left\langle R_k^{(\alpha, \beta)}, R_l^{(\alpha, \beta)} \right\rangle_J = \delta_{kl} h_k \quad (k, l = 0, 1, \dots; h_k > 0)$$

Dual Bernstein basis polynomials: important results

- Ciesielski, 1987 ($\alpha = \beta = 0$): definition, recurrence relation.
- Jüttler, 1998 ($\alpha = \beta = 0$): Bernstein-Bézier representation.
- Lewanowicz, W., 2006–2017 ($\alpha, \beta > -1$): recurrence relation, orthogonal expansion, "short" representation, constrained dual Bernstein polynomials, Bézier form, ...
- Rababah, Al-Natour, 2007: extended Jüttler's results to the case $\alpha, \beta > -1$.
- Chudy, W., from 2018 ($\alpha, \beta > -1$): differential-recurrence relations, differential equation, new recurrence relations (for dual Bernstein polynomials of the same degree), ...

Dual Bernstein basis polynomials: explicit formulas

- Recurrence relation

$$D_i^{n+1}(x; \alpha, \beta) = \left(1 - \frac{i}{n+1}\right) D_i^n(x; \alpha, \beta) + \frac{i}{n+1} D_{i-1}^n(x; \alpha, \beta) + \vartheta_i^n R_{n+1}^{(\alpha, \beta)}(x),$$

where ϑ_i^n – simple coefficient

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$$D_i^n(x; \alpha, \beta) = \sum_{k=0}^n C_k^{(\alpha, \beta)} Q_k(i; \beta, \alpha, n) R_k^{(\alpha, \beta)}(x),$$

where $Q_k(i; \beta, \alpha, n)$ are Hahn orthogonal polynomials, and $C_k^{(\alpha, \beta)}$ – simple coefficient

Dual Bernstein basis polynomials: explicit formulas

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- Short representations

$$D_i^n(x; \alpha, \beta) = A_{ni}^{(\alpha, \beta)} \sum_{k=0}^i \frac{(-i)_k}{(-n)_k} R_{n-k}^{(\alpha, \beta+k+1)}(x),$$

$$D_{n-i}^n(x; \alpha, \beta) = B_{ni}^{(\alpha, \beta)} \sum_{k=0}^i (-1)^k \frac{(-i)_k}{(-n)_k} R_{n-k}^{(\alpha+k+1, \beta)}(x),$$

where $A_{ni}^{(\alpha, \beta)}$, $B_{ni}^{(\alpha, \beta)}$ – simple coefficients

Dual Bernstein basis polynomials: newer results

- Representation in the basis $(1 - x)^j$ ($0 \leq j \leq n$)

$$D_i^n(x; \alpha, \beta) = \sum_{j=0}^n E_{nij}^{(\alpha, \beta)} {}_3F_2 \left(\begin{matrix} j - n, -i, 1 \\ -n, -n - \alpha \end{matrix} \middle| 1 \right) \cdot (1 - x)^j \quad (0 \leq i \leq n),$$

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- $F(i, j) := {}_3F_2 \left(\begin{matrix} j - n, -i, 1 \\ -n, -n - \alpha \end{matrix} \middle| 1 \right) \quad (i, j = 0, 1, \dots, n)$

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- The quantities $F(i, j)$ satisfy recurrence relation of the form

$$\begin{aligned} (i - n)(n - i + \alpha)F(i + 1, j) - (i + 1)(n + j - i + \alpha + 1)F(i, j) = \\ = -(n + 1)(n + \alpha + 1) \quad (0 \leq i, j \leq n) \end{aligned}$$

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- **Proof.** Zeilberger's algorithm

Dual Bernstein basis polynomials: newer results (cont.)

- Differential-recurrence relations

$$\begin{aligned} \left((1-x)\mathbf{D} - (n-i+\alpha+1)\mathbf{I} \right) \mathbf{D}_i^n(x; \alpha, \beta) &= \\ &= \frac{(i-n)(i+\beta+1)}{i+1} \mathbf{D}_{i+1}^n(x; \alpha, \beta) - \mathbf{u}_{ni}^{(\alpha, \beta)} \frac{n+\alpha+1}{i+1} \mathbf{R}_n^{(\alpha, \beta+1)}(x), \end{aligned}$$

where $0 \leq i \leq n$, $\mathbf{D} := \frac{d}{dx}$, \mathbf{I} – identity operator, a $\mathbf{u}_{ni}^{(\alpha, \beta)}$ – simple coefficient

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$$\begin{aligned} \left(x\mathbf{D} + (i+\beta+1)\mathbf{I} \right) \mathbf{D}_i^n(x; \alpha, \beta) &= \\ &= \frac{i(n-i+\alpha+1)}{n-i+1} \mathbf{D}_{i-1}^n(x; \alpha, \beta) + \mathbf{U}_{ni}^{(\alpha, \beta)} \frac{n+\beta+1}{n-i+1} \mathbf{R}_n^{(\alpha+1, \beta)}(x), \end{aligned}$$

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where $0 \leq i \leq n$, $\mathbf{D} := \frac{d}{dx}$, \mathbf{I} – identity operator, a $\mathbf{U}_{ni}^{(\alpha, \beta)}$ – simple coefficient.

- Proof.** Representation of \mathbf{D}_i^n in the basis $(1-x)^j$ + recurrence for $F(i, j)$ + symmetry:

$$\mathbf{D}_i^n(x; \alpha, \beta) = \mathbf{D}_{n-i}^n(1-x; \beta, \alpha)$$

Dual Bernstein basis polynomials: newer results (cont.)

- Third differential-recurrence relation

$$\begin{aligned} & \left(x(x-1)\mathbf{D}^2 + \frac{1}{2}(\alpha - \beta + (\sigma + 1)(2x - 1))\mathbf{D} \right) \mathbf{D}_i^n(x; \alpha, \beta) = \\ & = (i - n)(i + \beta + 1)\mathbf{D}_{i+1}^n(x; \alpha, \beta) + i(i - \alpha - n - 1)\mathbf{D}_{i-1}^n(x; \alpha, \beta) - \\ & \quad - (i(i - \alpha - n - 1) + (i - n)(i + \beta + 1))\mathbf{D}_i^n(x; \alpha, \beta), \end{aligned}$$

where $0, 1, \dots, n$

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where $0, 1, \dots, n$.

- **Proof.** Representation in the basis $\mathbf{R}_k^{(\alpha, \beta)}(x)$ + differential equation for $\mathbf{R}_k^{(\alpha, \beta)}(x)$ + difference equation for Hahn orthogonal polynomials $Q_k(x; \mathbf{a}, \mathbf{b}; \mathbf{N})$

Dual Bernstein basis polynomials: newer results (cont.)

- Fourth-order differential equation

$$\sum_{j=0}^4 w_j(x) \mathbf{D}^j D_i^n(x; \alpha, \beta) = 0 \quad (0 \leq i \leq n),$$

where $w_j(x)$ is a polynomial of degree j ($0 \leq j \leq 4$)

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- **Proof.** It follows from new differential-recurrence relations

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The *newer* results



- F. Chudy, P. W., *Differential-recurrence properties of dual Bernstein polynomials*, Applied Mathematics and Computation 338 (2018), 537–543

Dual Bernstein basis polynomials: newest results

- First-order non-homogeneous recurrence relation

$$\begin{aligned}
 (x - 1)(i + 1)D_i^n(x; \alpha, \beta) + x(n - i)D_{i+1}^n(x; \alpha, \beta) = \\
 C_{ni}^{(\alpha, \beta)} \left((n - i)(n + \alpha + 1)xR_n^{(\alpha, \beta+1)}(x) + \right. \\
 \left. + (i + 1)(n + \beta + 1)(1 - x)R_n^{(\alpha+1, \beta)}(x) \right),
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where $C_{ni}^{(\alpha, \beta)}$ – simple coefficient.

- **Proof.** Induction on n + properties of shifted Jacobi polynomials

Dual Bernstein basis polynomials: newest results (cont.)

- Third-order recurrence relation

$$\sum_{j=0}^3 v_j(n, i; x; \alpha, \beta) D_{i+j}^n(x; \alpha, \beta) = 0 \quad (0 \leq i \leq n-3),$$

where $v_j(n, i; x; \alpha, \beta)$ – simple cubic polynomial in i (and linear polynomial in x)
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Dual Bernstein basis polynomials: newest results (cont.)

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- **Proof.** It follows from new first-order non-homogeneous recurrence relation

Dual Bernstein basis polynomials: newest results (cont.)

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- F. Chudy, P. W., *Fast and accurate evaluation of dual Bernstein polynomials*, Numerical Algorithms 87 (2021), 1001–1015

Dual Bernstein basis polynomials: newest results (cont.)

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- **Problem.** Let us fix numbers: $n \in \mathbb{N}$, $x \in [-1, 1]$ and $\alpha, \beta > -1$. Consider the problem of computing the values

$$D_i^n(x; \alpha, \beta)$$

for all $i = 0, 1, \dots, n$

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- **Solution.**

First-order non-homogeneous recurrence relation $\Rightarrow O(n)$

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First-order non-homogeneous recurrence relation $\Rightarrow O(n)$.

Previous results $\Rightarrow O(n^2)$

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- Tests: veery good numerical properties :-)

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- (Classical) Dual Bernstein polynomials:
 - three differential-recurrence relations,
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- (Classical) Dual Bernstein polynomials:
 - three differential-recurrence relations,
 - fourth-order differential equation,
 - first-order non-homogeneous recurrence relation,
 - third-order recurrence relation.
- For a given x , α , β , the values of all $n + 1$ (classical) dual Bernstein polynomials $D_k^n(x; \alpha, \beta)$ ($0 \leq k \leq n$) can be computed with the complexity $O(n)$

